

# Approximation Algorithms

## problem set #1 – due Wednesday, April 19

**Problem 1 (Two-Dimensional Bin Packing).** Consider the following problem. We are given a collection of items  $1, \dots, n$ . Each item  $i$  has a weight  $w_i$  and size  $s_i$ . Our goal is to pack all items in the minimum number of identical bins so that the total weight of items in every bin is at most  $W$ , and the total size is at most  $S$ .

Give a constant factor approximation algorithm for the problem.

### Problem 2 (Max Cut)

- Give a 2-approximation algorithm for the Max Cut problem. In Max Cut, we are given a graph  $G = (V, E)$ , and our goal is to partition the set of vertices  $V$  into two sets  $L$  and  $R$  to maximize the size of the cut between  $L$  and  $R$ . The size of the cut between  $L$  and  $R$  equals the number of edges with one endpoint in  $L$  and the other in  $R$ .
- Give a 4-approximation algorithm for Directed Max Cut. Given a *directed* graph  $G = (V, A)$ , partition the set of vertices  $V$  into two sets  $L$  and  $R$  to maximize the size of the *directed* cut between  $L$  and  $R$ . The size of the *directed* cut between  $L$  and  $R$  equals the number of arcs  $(u, v) \in A$  with  $u \in L$  and  $v \in R$ .

**Problem 3.** Suppose we want to schedule  $n$  jobs  $1, \dots, n$  on a single machine. Each job  $j$  has a processing time  $p_j$  and weight  $w_j$ . The completion time  $C_j$  of job  $j$  is the time when  $j$  is completed. Give an exact algorithm for finding the schedule with the minimum weighted completion time defined as follows:

$$\text{weighted completed time} = \sum w_j C_j$$

**Problem 4.** Give an example of a set of points  $X$  in  $\mathbb{R}^d$  and an initial set of centers (seed for the algorithm)  $c_1, \dots, c_k$  in  $\mathbb{R}^d$  for which Lloyd's algorithm returns a suboptimal solution. What is the approximation of the algorithm in this case?